

When do quasisymmetric functions know that trees are different?

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Joint work with:
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Discrete Mathematics Seminar, Xiamen University
30 November 2022



Slides and paper available from
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- ▶ Chromatic (quasi)symmetric functions and the motivating conjectures
- ▶ Converting to a poset question; more conjectures
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The chromatic polynomial

George Birkhoff, 1912

Graph $G = (V, E)$

Colouring/Coloring: a map $\kappa : V \rightarrow \{1, 2, 3, \dots\}$

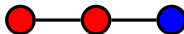
Proper coloring: adjacent vertices
get different colors.



Proper



Not Proper



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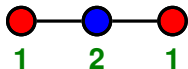
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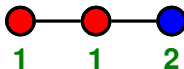
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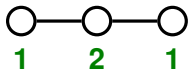
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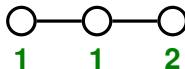
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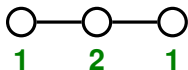
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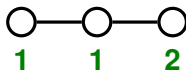
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Chromatic polynomial: $\chi_G(k)$ is the number of proper colorings of G when k colors are available.

Example.

$$\chi_G(k) = k(k - 1)(k - 1)$$

The chromatic symmetric function

Richard Stanley, 1995

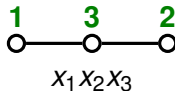
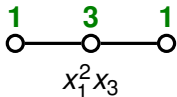


Graph $G = (V, E)$

$$V = \{v_1, v_2, \dots, v_n\}$$

To a proper coloring κ , we associate the monomial in commuting variables x_1, x_2, \dots

$$x_{\kappa(v_1)} x_{\kappa(v_2)} \cdots x_{\kappa(v_n)}$$



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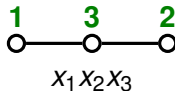
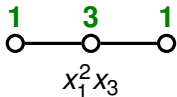


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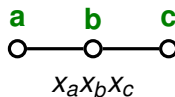
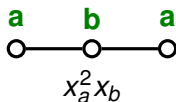
$$X_G(x_1, x_2, \dots) = X_G(\mathbf{x}) = \sum_{\text{proper } \kappa} x_{\kappa(v_1)} x_{\kappa(v_2)} \cdots x_{\kappa(v_n)}.$$

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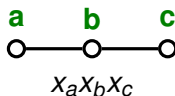
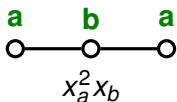


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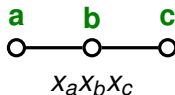
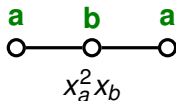
$$X_G(\mathbf{x}) = \sum_{a \neq b} x_a^2 x_b + 6 \sum_{a < b < c} x_a x_b x_c$$
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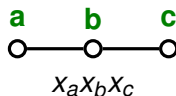
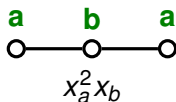
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- ▶ $X_G(\mathbf{x})$ is a symmetric function (invariant when you permute the colors/variables)
- ▶ Setting $x_i = 1$ for $1 \leq i \leq k$ and $x_i = 0$ otherwise yields $\chi_G(k)$.
e.g. $k(k-1) + 6\binom{k}{3} = k(k-1)^2$.

Can $X_G(\mathbf{x})$ distinguish graphs?

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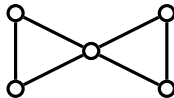
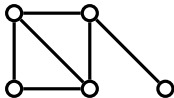
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$X_G(\mathbf{x})$ distinguishes **trees**. In other words, if T and U are non-isomorphic trees, then $X_T(\mathbf{x}) \neq X_U(\mathbf{x})$.



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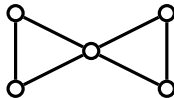
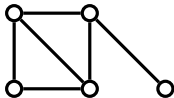
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[Aliste-Prieto, Crew, de Mier, Fougere, Heil, Ji, Loeb, Martin, Morin, Orellana, Scott, Smith, Sereni, Spirkl, Tian, Wagner, Zamora, ...]

Remark. Steph van Willigenburg: another famous $X_G(\mathbf{x})$ conjecture.

A little bit of (quasi)symmetric functions

$x^2y + y^2x + x^2z + z^2x + y^2z + z^2y$ is a **symmetric polynomial** in $\{x, y, z\}$ because it doesn't change when you permute the variables.

$\sum_{a \neq b} x_a^2 x_b = x_1^2 x_2 + x_2^2 x_1 + x_1^2 x_3 + \dots$ is a **symmetric function** in \mathbf{x} .

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Now consider $\sum_{a < b} x_a x_b^2 = x_1 x_2^2 + x_1 x_3^2 + x_2 x_3^2 + x_1 x_4^2 + \dots$.

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For a composition $\alpha = (\alpha_1, \alpha_2, \dots, \alpha_k)$ the **monomial quasisymmetric function** is:

$$M_\alpha = \sum_{a_1 < a_2 < \dots < a_k} x_{a_1}^{\alpha_1} x_{a_2}^{\alpha_2} \dots x_{a_k}^{\alpha_k}.$$

The span of the M_α is the vector space **QSym** of **quasisymmetric functions**.

In fact,

- ▶ the M_α form a basis for $QSym$;
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A more important basis for us is Gessel's **fundamental quasisymmetric functions**:

$$F_\alpha = \sum_{\beta \text{ refines } \alpha} M_\beta.$$

Example.

$$F_{32} = M_{32} + M_{212} + M_{122} + M_{1112} + M_{311} + M_{2111} + M_{1211} + M_{11111}.$$

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$QSym$ is a star of 21st century algebraic combinatorics.

The chromatic **quasi**symmetric function

John Shareshian & Michelle Wachs, 2014; Brittney Ellzey, 2017.

Directed graph $\vec{G} = (V, E)$.

Ascent of proper coloring κ : directed edge $u \rightarrow v$ with $\kappa(u) < \kappa(v)$

asc(κ): the number of ascents of κ .

Example. Colors $a < b < c$



$\kappa(V_1)$	$\kappa(V_2)$	$\kappa(V_3)$	asc(κ)
a	b	c	1
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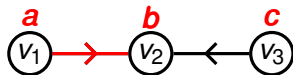
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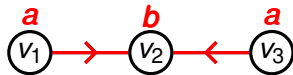
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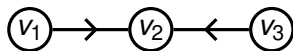
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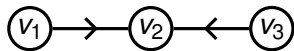
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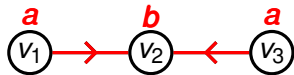
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By setting $t = 1$, we see that $X_{\vec{G}}(\mathbf{x}, t)$ contains more information than $X_G(\mathbf{x})$.

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Statement 3.

$X_{\vec{G}}(\mathbf{x}, t)$ distinguishes directed graphs.

i.e. if \vec{G} and \vec{H} are not isomorphic, then $X_{\vec{G}}(\mathbf{x}, t) \neq X_{\vec{H}}(\mathbf{x}, t)$.



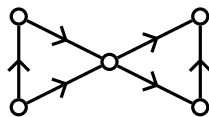
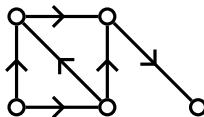
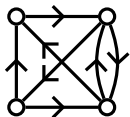
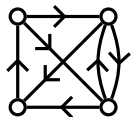
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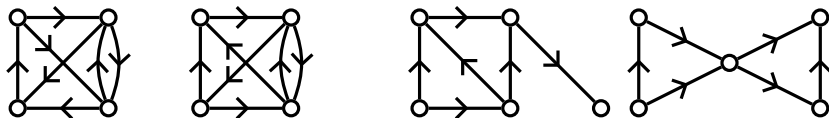
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Statement 4.

$X_{\vec{G}}(\mathbf{x}, t)$ distinguishes directed **trees**. In other words, if \vec{T} and \vec{U} are non-isomorphic directed trees, then $X_{\vec{T}}(\mathbf{x}, t) \neq X_{\vec{U}}(\mathbf{x}, t)$.



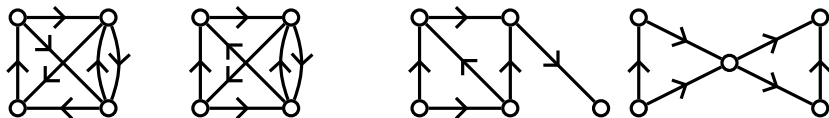
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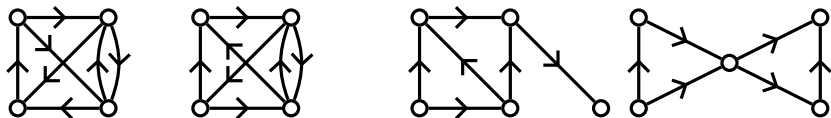
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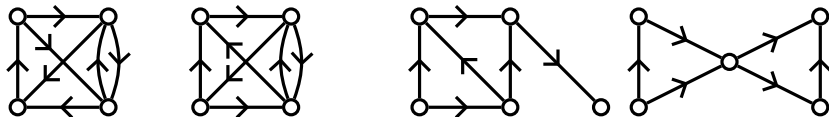
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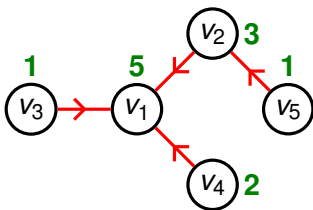
This conjecture was our original goal. **Strategy: translate to posets.**

$$X_{\vec{G}}(\mathbf{x}, t) = \sum_{\text{proper } \kappa} t^{\text{asc}(\kappa)} X_{\kappa(v_1)} X_{\kappa(v_2)} \cdots X_{\kappa(v_n)}.$$

Want to show: $X_{\vec{T}}(\mathbf{x}, t) \neq X_{\vec{U}}(\mathbf{x}, t)$.

Key insight:

- ▶ Look at the coefficient of the highest power of t .
- ▶ It's enough to show these coefficients are different for T and U .
- ▶ So just look at colorings where all edges are ascents

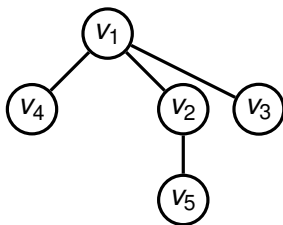
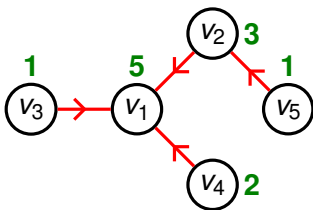


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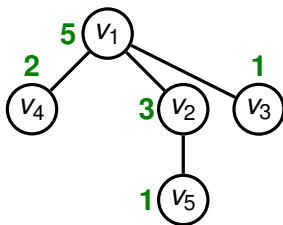
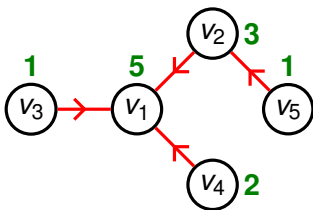


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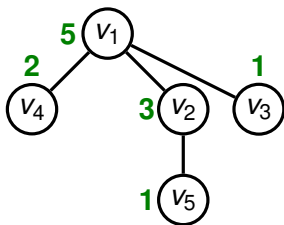
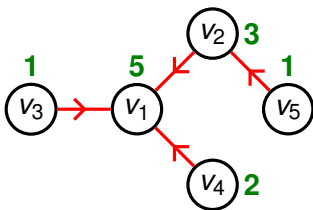


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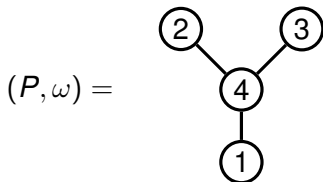
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- ▶ The corresponding coloring is a **strict P -partition**.



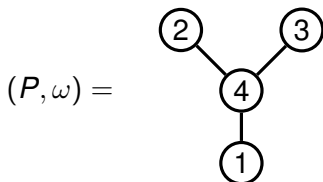
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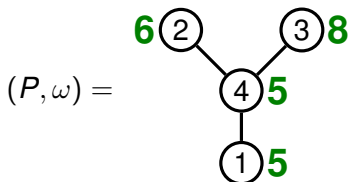


Key definition. A (P, ω) -**partition** is a map f from P to the positive integers satisfying:

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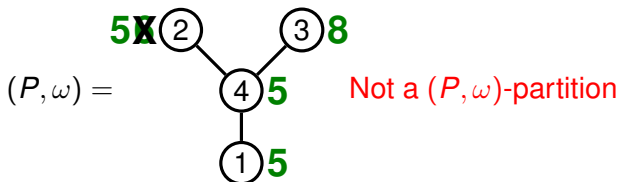


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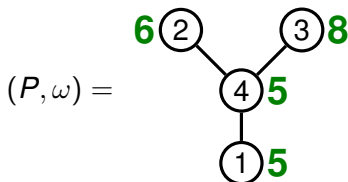


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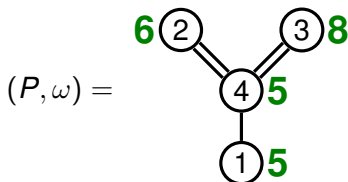


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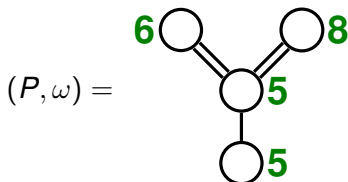
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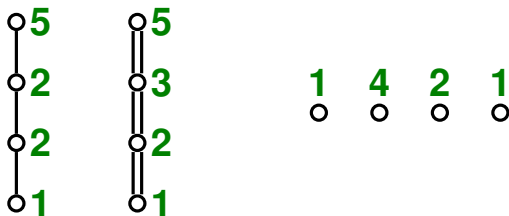


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We use double edges to denote the strictness conditions and then we can (usually) ignore the underlying labeling.

Motivating examples for (P, ω) -partitions

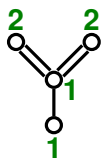


- ▶ (P, ω) chain with all weak edges: get a partition
- ▶ (P, ω) chain with all strict edges: get a partition with distinct parts
- ▶ (P, ω) is an antichain: get a composition

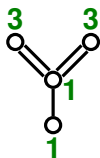
General (P, ω) -partitions interpolate between these classical objects.

The (P, ω) -partition enumerator

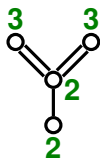
Example. Restrict to $f(p) \in \{1, 2, 3\}$.



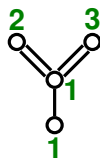
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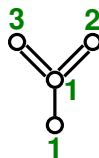
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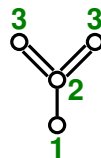
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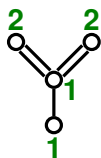
$$K_{(P, \omega)}(x_1, x_2, x_3) = x_1^2 x_2^2 + x_1^2 x_3^2 + x_2^2 x_3^2 + 2x_1^2 x_2 x_3 + x_1 x_2 x_3^2.$$

In general, the (P, ω) -partition enumerator is by given by:

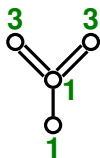
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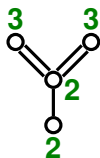
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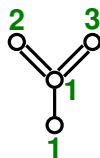
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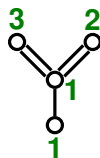
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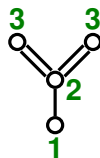
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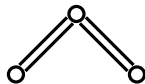
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Seem familiar?

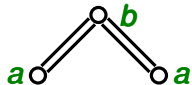
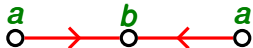
From $X_{\vec{G}}(\mathbf{x}, t)$ to $K_{(P, \omega)}$



From $X_{\vec{G}}(\mathbf{x}, t)$ to $K_{(P, \omega)}$

colorings of \vec{G} will all ascents \longleftrightarrow strict P -partitions

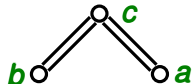
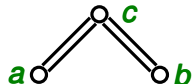
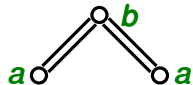
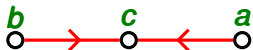
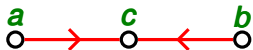
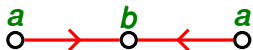
$$a < b < c$$



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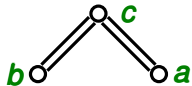
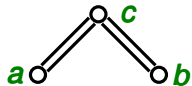
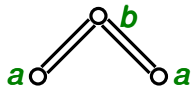
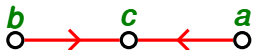
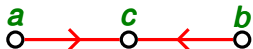
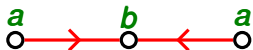
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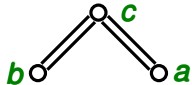
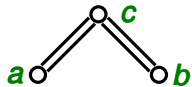
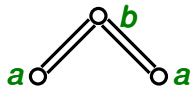
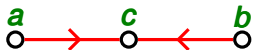
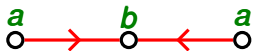


$$\text{coefficient of } t^2 \text{ in } X_{\vec{G}}(\mathbf{x}, t) = M_{21} + 2M_{111} = K_P^<(\mathbf{x}).$$

From $X_{\vec{G}}(\mathbf{x}, t)$ to $K_{(P, \omega)}$

colorings of \vec{G} will all ascents \longleftrightarrow strict P -partitions

$a < b < c$



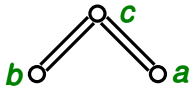
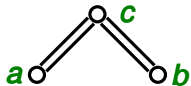
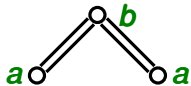
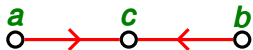
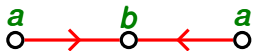
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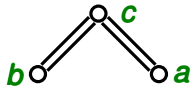
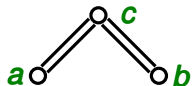
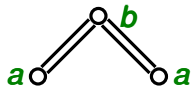
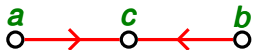
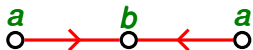
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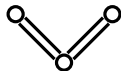
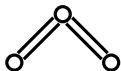
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[Browning, Féray, Hasebe, Hopkins, Kelly, Liu, M., Tsujie, Ward, Weselcouch]

Can $K_{(P,\omega)}(\mathbf{x})$ distinguish posets?



Statement 5.

$K_P^{\leq}(\mathbf{x})$ distinguishes posets that are trees.

i.e. if tree posets P and Q are not isomorphic, then $K_P^{\leq}(\mathbf{x}) \neq K_Q^{\leq}(\mathbf{x})$.



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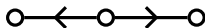
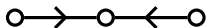
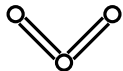


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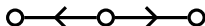
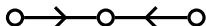
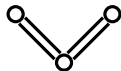
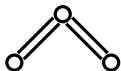
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Statement 6. (mix strict and weak edges)

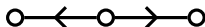
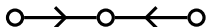
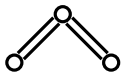
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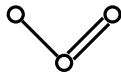
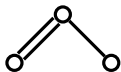
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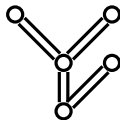
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Statement 7.

$K_P^<(\mathbf{x})$ distinguishes posets that are **rooted** trees.

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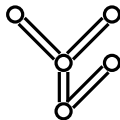
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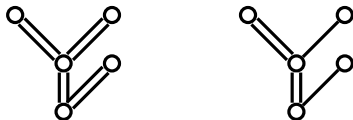
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We'd like to allow a mixture of strict and weak edges

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Conjecture 4 [Aval, Djenabou, M., 2022].

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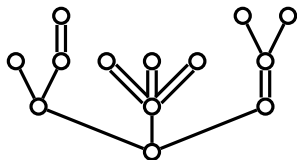
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Our main contribution sits between Theorem 1 and Conjecture 4.

Fair trees and a generalization

Definition. A labeled poset that is a tree is said to be a **fair tree** if for each vertex, its outgoing edges up to its children are either all strict or all weak.

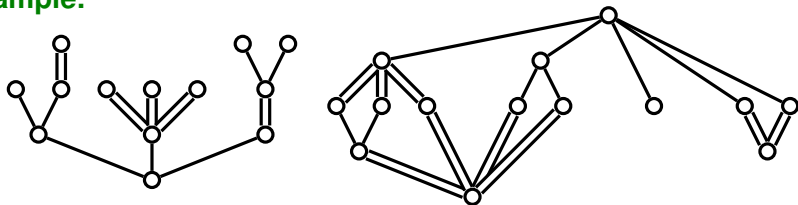
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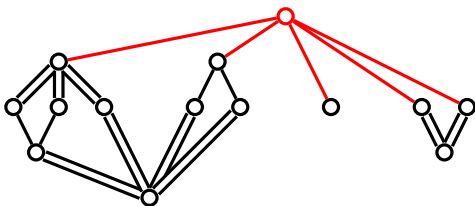
Definition. More generally, we define the set \mathcal{C} of labeled posets recursively by:

1. the one-element labeled poset $[1]$ is in \mathcal{C} ;
2. \mathcal{C} is closed under disjoint unions $(P, \omega) \sqcup (Q, \omega')$ is in \mathcal{C} ;
3. \mathcal{C} is closed under the ordinal sums $(P, \omega) \uparrow [1]$ and $(P, \omega) \uparrow\uparrow [1]$;
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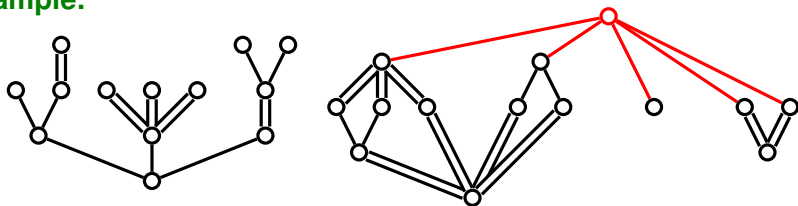
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Our main theorem

Statement 9.

$K_{(P,\omega)}(\mathbf{x})$ distinguishes elements of \mathcal{C} , so in particular fair trees;
i.e. if (P, ω) and (Q, τ) are in \mathcal{C} and not isomorphic, then
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Crux of the proof:

Proposition 1 [Aval, Djenabou, M., 2022]

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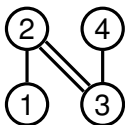
Irreducibility is also the crux for

- ▶ Hasebe & Tsujie;
- ▶ Ricki Ini Liu & Michael Weselcouch ($K_P^{\leq}(\mathbf{x})$ distinguishes series-parallel posets; needs irreducibility for general P with all strict edges, 2020).

Main tool in this research area

Stanley, 1971 and Ira Gessel, 1984:
 $K_{(P,\omega)}(\mathbf{x})$ expands beautifully in F -basis.

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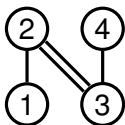


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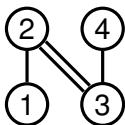


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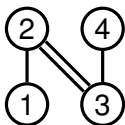
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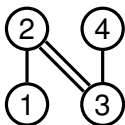
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$$K_{(P,\omega)} = 2F_{22} + F_{31} + F_{13} + F_{121}.$$

Theorem [Gessel & Stanley]. For a labeled poset (P, ω) ,

$$K_{(P,\omega)} = \sum_{\pi \in \mathcal{L}(P,\omega)} F_{\text{comp}(\pi)}.$$

Some final conjectures

Recall Stanley's

Famous Conjecture 1. $X_G(\mathbf{x})$ distinguishes **trees**. In other words, if T and U are non-isomorphic trees, then $X_T(\mathbf{x}) \neq X_U(\mathbf{x})$.

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Surprising Conjecture 5 [Nick Loehr & Greg Warrington, 2022].

$X_G(1, q, q^2, \dots, q^{n-1})$ distinguishes trees with n vertices, i.e. if T and U are non-isomorphic trees with n vertices, then

$$X_T(1, q, q^2, \dots, q^{n-1}) \neq X_U(1, q, q^2, \dots, q^{n-1}).$$

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Remark. This specialization has a nice interpretation for $K_{(P, \omega)}$: if

$$K_{(P, \omega)}(1, q, q^2, \dots, q^{k-1}) = \sum_{N \geq 0} a(N) q^N,$$

then we see that $a(N)$ counts the number of (P, ω) -partitions $f: P \rightarrow \{0, \dots, k-1\}$ of N .

Some final conjectures

Recall **Conjecture 3**. $K_P^{\leq}(\mathbf{x})$ distinguishes posets that are trees, i.e. if tree posets P and Q are not isomorphic, then $K_P^{\leq}(\mathbf{x}) \neq K_Q^{\leq}(\mathbf{x})$.

Conjecture 6 [Aval, Djenabou, M., 2022].

$K_P^{\leq}(1, q, q^2, \dots, q^{n-1})$ distinguishes tree posets with n elements, i.e. if T and U are non-isomorphic trees with n vertices, then

$$K_P^{\leq}(1, q, q^2, \dots, q^{n-1}) \neq K_U^{\leq}(1, q, q^2, \dots, q^{n-1}).$$

Remark. This specialization has a nice interpretation for $K_{(P, \omega)}$: if

$$K_{(P, \omega)}(1, q, q^2, \dots, q^{k-1}) = \sum_{N \geq 0} a(N) q^N,$$

then we see that $a(N)$ counts the number of (P, ω) -partitions $f: P \rightarrow \{0, \dots, k-1\}$ of N .

Thanks for your attention! 謝謝